

# EE 332 Real Time Systems

## Midterm Examination Solution

Wednesday February 12, 2003

### General:

- Two hours (2:30pm to 4:30pm)
- Open book and open notes

### 1. 15 Marks (3 marks each)

- a) Briefly describe the difference between call-by-value and call-by-reference and give an example of each.

Call-by-value passes the value of a parameter to a function. Thus the function can only “read” the value of the parameter. Call-by-value parameters require one indirect access.

```
int add( int a, int b )
{
    a = a + b;
    return a;
}
```

Call-by-reference passes the address (or reference) of a parameter to a function. Thus the function can both “read” from and “write” to the parameter. Call-by-reference parameters require two indirect accesses.

```
int add( int* a, int* b )
{
    *a = *a + *b;
    return *a;
}
```

- b) Some processors provide an instruction that allows the processor to go to sleep until an interrupt occurs. Why is this desirable and useful?

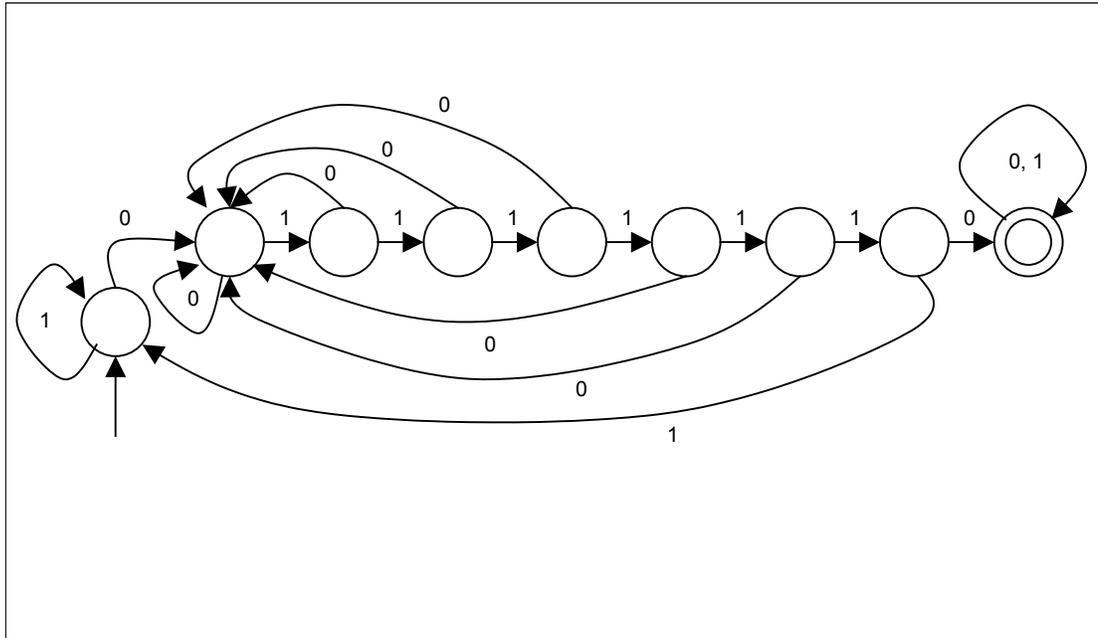
The ability for a processor “to go to sleep until an interrupt occurs” is very useful for low-power consumption (e.g., battery powered).

- c) Is the following routine thread-safe? Provide a “yes”, “no”, or “cannot determine” answer and give a brief justification for your answer:

```
int    GlobalCount;
void IncGlobalCount( void )
{
    GlobalCount++;
}
```

The correct answer is “cannot determine”. If the global variable can be incremented “atomically” (i.e., in a single processor instruction), it cannot be interrupted and the routine would be thread-safe. If however the global variable cannot be incremented atomically (i.e., requires more than one processor instruction), it could be interrupted and the routine would not be thread-safe. Depends on the particular processor architecture.

- d) Show the Moore Finite State Automata for a system that scans for the bit sequence 01111110 in a string of bits (the sequence is allowed to appear anywhere in the bit string).



- e) Why is it important for a real-time system to have a watchdog timer? Is a timer interrupt routine an appropriate point in the software to reset the watchdog timer, why or why not?

It is important for a real-time system to have a watchdog timer to reset the processor if it gets into an unintended state (e.g., infinite loop). This is especially important for embedded and unattended applications. A timer interrupt routine is NOT an appropriate point to reset the watchdog timer as it just proves that interrupts are working, it does not guard against improper operation in the foreground of the software (e.g., infinite loop).

## 2. 20 Marks (10 + 10)

Fibonacci numbers follow the sequence 1, 1, 2, 3, 5, 8, ... Except for the first two numbers, each number is the sum of the preceding two numbers. A recursive routine for determining a particular number in the sequence is given by:

```
int fibonacci( int number )
{
    if( number < 2 )
    {
        return 1;
    }
    else
    {
        return( fibonacci( number - 1 ) + fibonacci( number - 2 ) );
    }
}
```

- a) Show how the fibonacci recursive algorithm works by showing the values of the parameters to each call of fibonacci( ), how many times fibonacci( ) is called, and the return value of each fibonacci( ) when you call “fibonacci( 5)”.

```

fibonacci( 5 )
  fibonacci( 4 )
    fibonacci( 3 )
      fibonacci( 2 )
        fibonacci( 1 )
          return 1
        fibonacci( 0 )
          return 1
        return 2
      fibonacci( 1 )
        return 1
      return 3
    fibonacci( 2 )
      fibonacci( 1 )
        return 1
      fibonacci( 0 )
        return 1
      return 2
    return 5
  fibonacci( 3 )
    fibonacci( 2 )
      fibonacci( 1 )
        return 1
      fibonacci( 0 )
        return 1
      return 2
    fibonacci( 1 )
      return 1
    return 3
  return 8

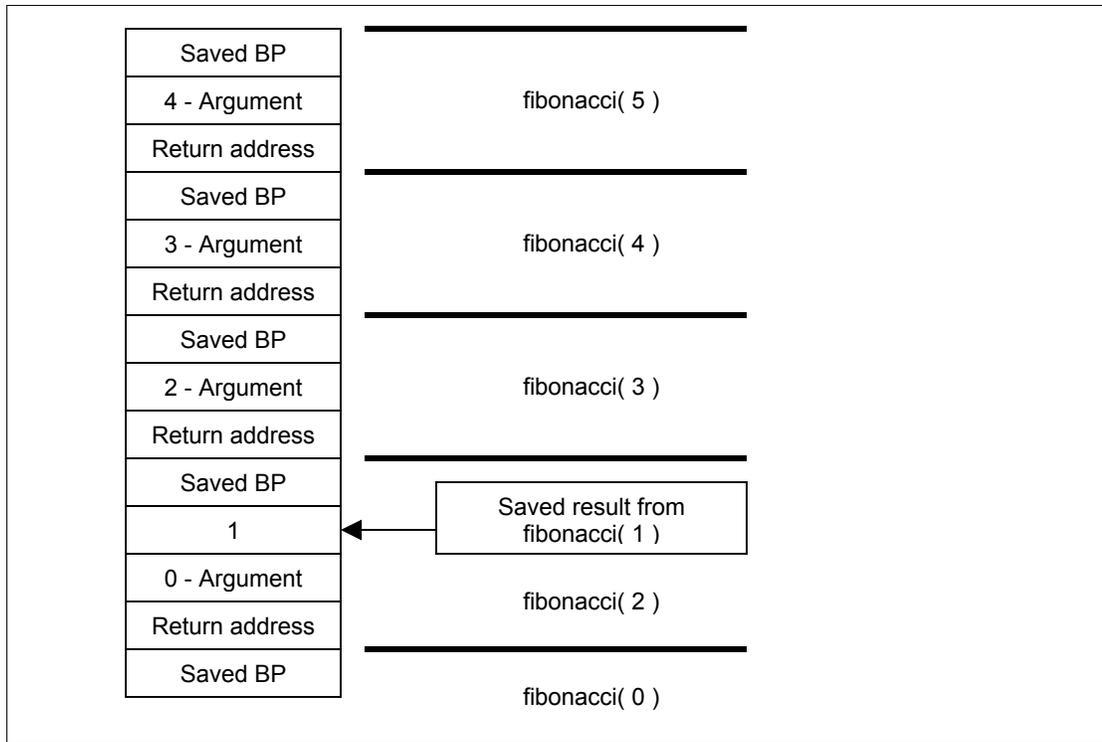
```

- b) Given the 8086 assembler code shown from the fibonacci C code, show the maximum length call stack when you call “fibonacci( 5)”. Indicate the value and what is represented by each word of the stack.

```

_fibonacci
    push    bp
    mov     bp,sp
    cmp     word ptr [bp+4],2
    jge     short @1@86
    mov     ax,1
    jmp     short @1@114
@1@86:
    mov     ax,word ptr [bp+4]
    dec     ax
    push    ax
    call    near ptr _fibonacci
    pop     cx
    push    ax
    mov     ax,word ptr [bp+4]
    add     ax,-2
    push    ax
    call    near ptr _fibonacci
    pop     cx
    pop     dx
    add     dx,ax
    mov     ax,dx
@1@114:
    pop     bp
    ret

```



### 3. 25 Marks (10 + 10 + 5)

Consider a hypothetical microprocessor called the CPU332. The CPU332 has the following properties:

- 16-bit microprocessor (16-bit data bus, registers, etc.)
- One 16-bit accumulator ACC
- Four 16-bit general purpose registers R0, R1, R2, R3
- A 16-bit stack pointer SP
- A 16-bit flag word FLAGS. Bit 0 (Least Significant Bit) is the “Interrupt Enable” bit, a value of “1” enables interrupts, a value of “0” disables interrupts.
- A “push” instruction can push any register (ACC, R0-R3, FLAGS) onto the stack. It does this by first writing the value and then incrementing the value of SP by two.
- A “pop” instruction can pop any register (ACC, R0-R3, FLAGS) from the stack. It does this by first decrementing the value of SP by two and then reading the value.

- a) Show the “pseudo assembler” for a “Yield( )” function for the CPU332. The pseudo assembler should show the order of pushes, pops, and saving and restoring of stack pointers. Indicate which registers are pushed and in what order. The code to determine the next task to run can be glossed over (but you should still indicate where in your Yield( ) function this occurs). Note that this Yield( ) function should be usable for either a cooperatively or preemptive multitasking system.

Yield:

```

push  FLAGS
push  ACC
push  R0
push  R1
push  R2
push  R3

mov   SP, SaveStack[ ActiveTask ]    ; Save SP

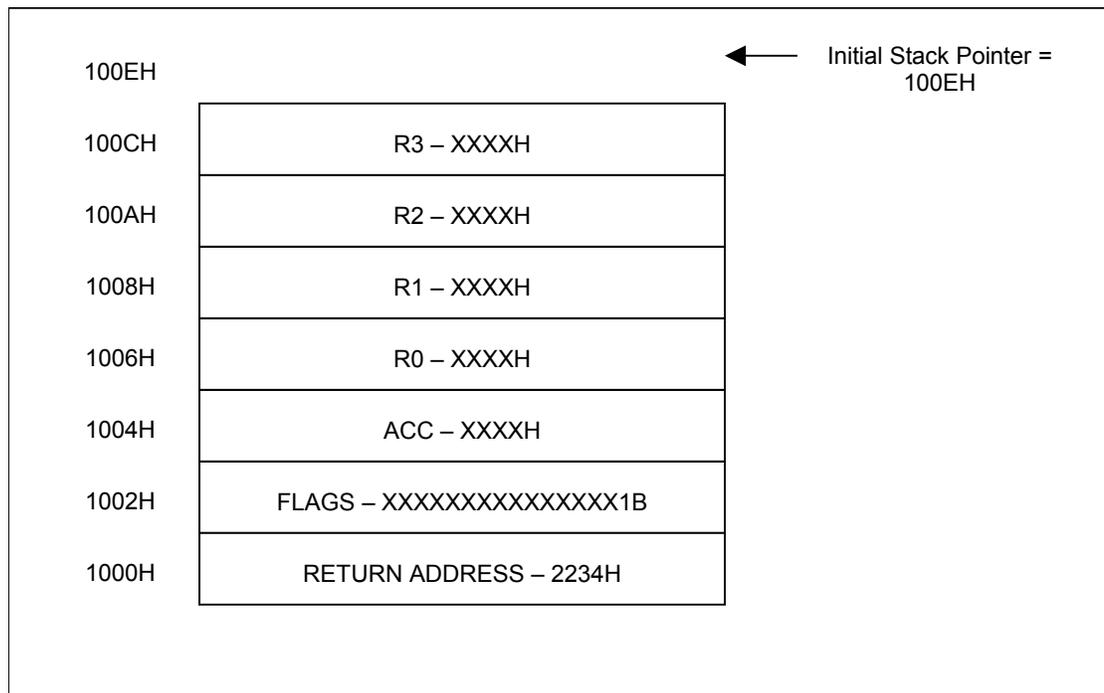
;; Code here to determine next task to run.

mov   SaveStack[ ActiveTask ], SP    ; Restore SP

pop   R3
pop   R2
pop   R1
pop   R0
pop   ACC
pop   FLAGS
ret

```

- b) For your Yield( ) function from (a), show the organization, addresses and values of the initial stack for a Task whose entry address is 2234H and whose stack space starts at 1000H. Remember to show required initial values for any values on the stack (or XXXX for don't cares). What is the value of the initial stack pointer?



c) If the interface of created tasks is required to be:

```
void Task( short taskNum, char* taskName );
```

`taskNum` is a 16-bit integer indicating the number of this task.

`taskName` is a 16-bit pointer to a string which is the name of this task.

Show how the initial stack is now different from part (b) as you now need to pass these parameters to the initial entry of `Task`. Do not worry about the value for these two parameters but do show where on the initial stack they reside (Hint: Remember C pushes parameters from right to left, rightmost parameter is pushed first, leftmost parameter is pushed last). What is the value of the initial stack pointer now?

